

# A Bottom-Up Deterministic Weighted Tree Automaton for the $n$ -Gram Yield Function

Matthias Büchse   Tobias Denking   Heiko Vogler

Institute of Theoretical Computer Science  
Technische Universität Dresden

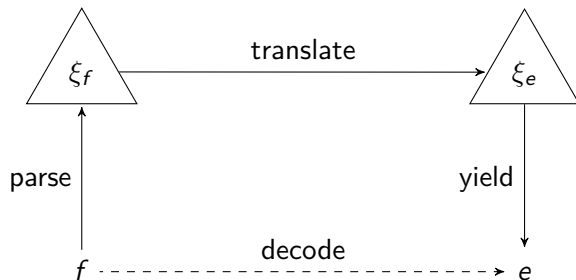
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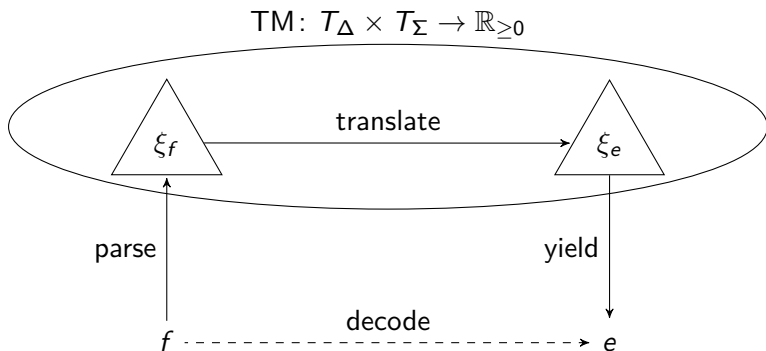
# Motivation: Statistical Machine Translation

$f$  ----- decode ----->  $e$

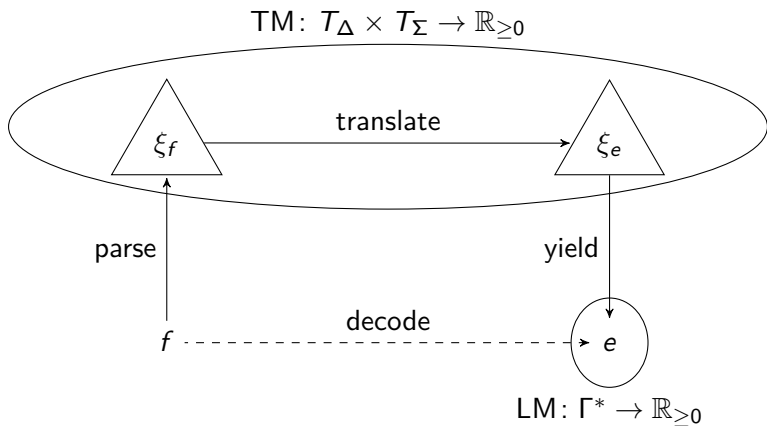
# Motivation: Statistical Machine Translation



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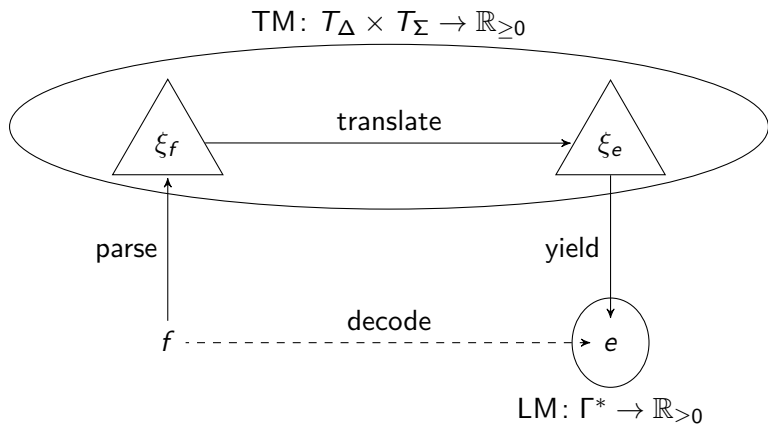
# Motivation: Statistical Machine Translation



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$$\Gamma = \Sigma^{(0)}$$

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$$\hat{e} = \pi_{\text{out}} \left( \operatorname{argmax}_d \left( (f \triangleleft TM) \triangleright LM \right) (d) \right)$$

---

$$\Gamma = \Sigma^{(0)}$$

# Generalized $n$ -Gram Models

Let  $n \in \mathbb{N} \setminus \{0\}$ .

A *generalized  $n$ -gram model* is a tuple  $N = (\Gamma, \mu_1, \dots, \mu_n)$  with

- the alphabet  $\Gamma$
- a mapping  $\mu_i: \Gamma^i \rightarrow \mathbb{R}_{\geq 0}$  for every  $i \in \{1, \dots, n\}$



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The semantics of  $N$  is  $\llbracket N \rrbracket: \Gamma^* \rightarrow \mathbb{R}_{\geq 0} \cup \{\perp\}$  where

$$\llbracket N \rrbracket(w) = \begin{cases} \prod_{w' \text{ } n\text{-gram of } w} \mu_n(w') & \text{if } |w| \geq n \\ \mu_{|w|}(w) & \text{if } 1 \leq |w| < n \\ \perp & \text{if } |w| = 0 \end{cases}$$

## Generalized $n$ -Gram Models (Example)

Let  $n = 3$ ,  $\Gamma = \{\text{Garcia, y, tres, asociados, .}\}$ , and

$\mu_3$ : Garcia y tres	$\mapsto 1/5$	$\mu_2$ : Garcia .	$\mapsto 1/2$	...
y tres asociados	$\mapsto 1/3$	...		
tres asociados .	$\mapsto 1/4$			
...				

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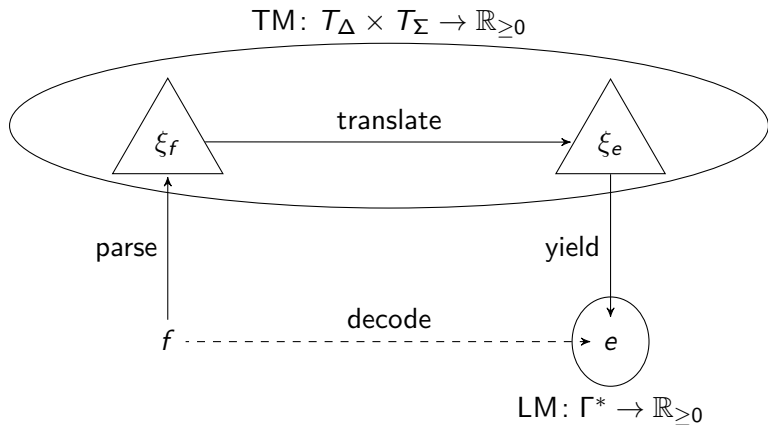
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$$\llbracket N \rrbracket(\varepsilon) = \perp$$

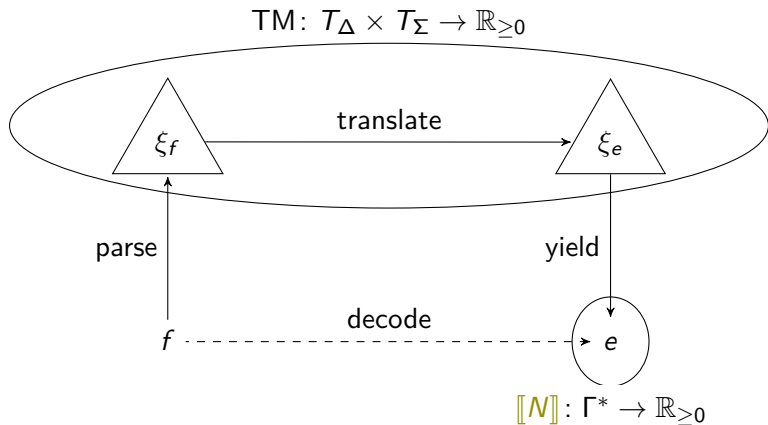


# Recall: Statistical Machine Translation



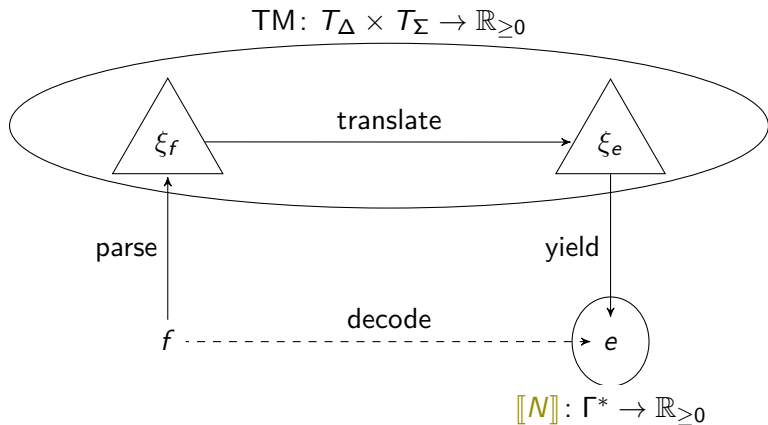
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# Recall: Statistical Machine Translation



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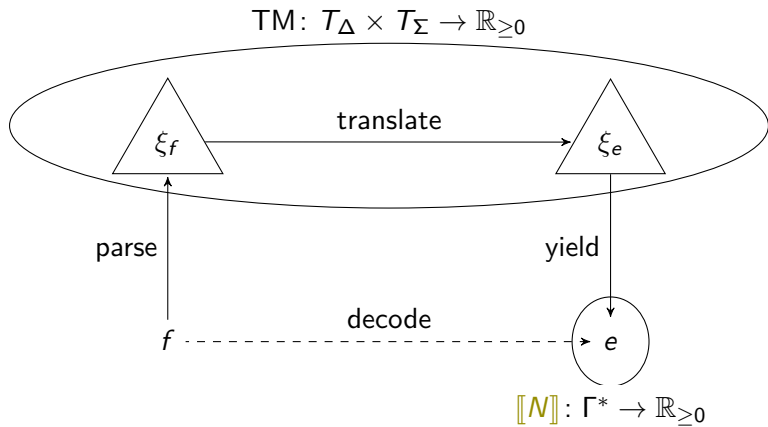
# Recall: Statistical Machine Translation



$$\hat{e} = \pi_{\text{out}} \left( \underset{T_{\Delta} \times T_{\Sigma} \rightarrow \mathbb{R}_{\geq 0}}{\text{argmax}_d} \left( (f \triangleleft TM) \triangleright [M] \right) (d) \right)$$

$T_{\Delta} \times T_{\Sigma} \rightarrow \mathbb{R}_{\geq 0}$        $\Gamma^* \rightarrow \mathbb{R}_{\geq 0}$

# Recall: Statistical Machine Translation



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$\Gamma^* \rightarrow \mathbb{R}_{\geq 0}$

lift to trees  
 $\text{yield}^{-1}([N])$

# Lifting using Bar-Hillel, Perles, Shamir (BHPS)

- for every  $\mathcal{L} \in \text{CF}$  and  $R \in \text{REG}$ :

$$\mathcal{L} \cap R \in \text{CF}$$

[Bar-Hillel, Perles, Shamir, 1961]

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- for every  $\mathcal{L} \in \text{WRT}$  and  $R \in \text{WREG}$ :  $\mathcal{L} \odot \text{yield}^{-1}(R) \in \text{WRT}$   
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- in our case:  $\mathcal{L}$  is the characteristic mapping of  $T_\Sigma$  and  $R = \llbracket N \rrbracket$   
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  - the constructed WTA is *not* bottom-up deterministic
  - *but* for every tree, only one run is relevant

---

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# Intermezzo: Weighted Tree Automata (WTA)

A *weighted tree automaton* is a tuple  $\mathcal{A} = (Q, \Sigma, \delta, F)$  with

- $Q$  finite set (states)
- $\Sigma$  ranked alphabet (input symbols)
- $\delta = (\delta_k \mid k \geq 0)$  with  $\delta_k: Q^k \times \Sigma^{(k)} \times Q \rightarrow \mathbb{R}_{\geq 0}$  (transition function)
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The *semantics of  $\mathcal{A}$*  is  $\llbracket \mathcal{A} \rrbracket: T_\Sigma \rightarrow \mathbb{R}_{\geq 0}$  where

$$\llbracket \mathcal{A} \rrbracket(\xi) = \sum_{\substack{\kappa \text{ run on } \xi \\ \kappa(\varepsilon) \in F}} \text{wt}(\kappa).$$

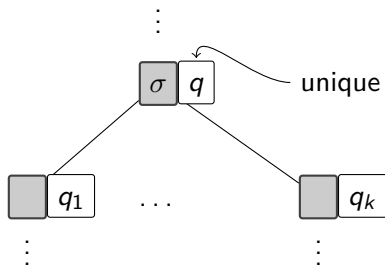
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$\sigma \dots$  label in  $\xi$  at  $\rho$

$k \dots$  rank of  $\sigma$

# Intermezzo: Weighted Tree Automata (WTA)

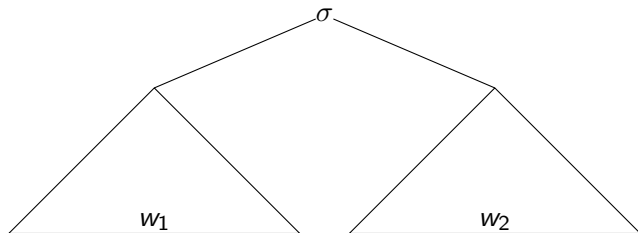
$\mathcal{A}$  is *bottom-up deterministic*:



$$\delta_k(q_1, \dots, q_k, \sigma, q) > 0$$

# Generalized $n$ -Gram WTA (Idea)

$$n = 3$$



$$\llbracket N \rrbracket(w_1 w_2) =$$

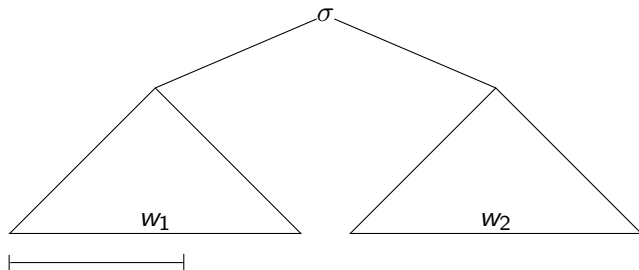
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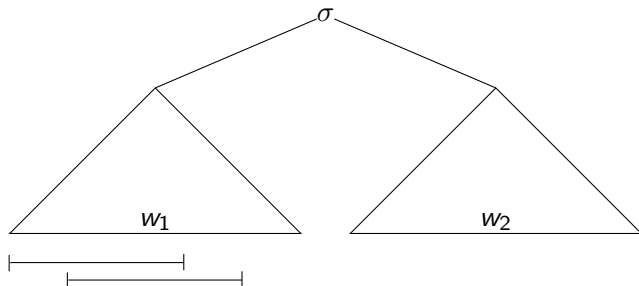
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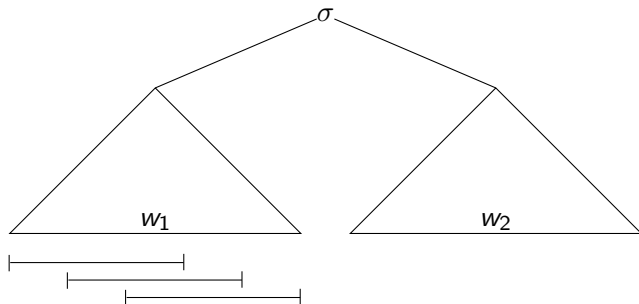
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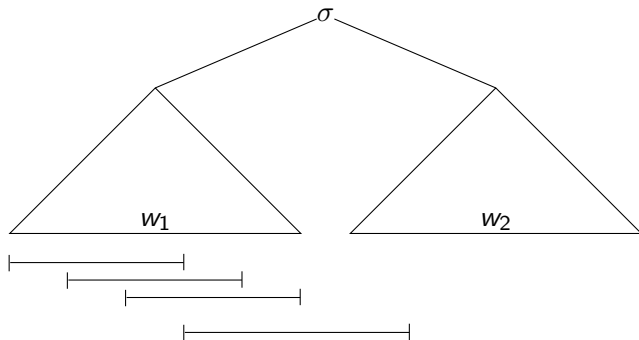
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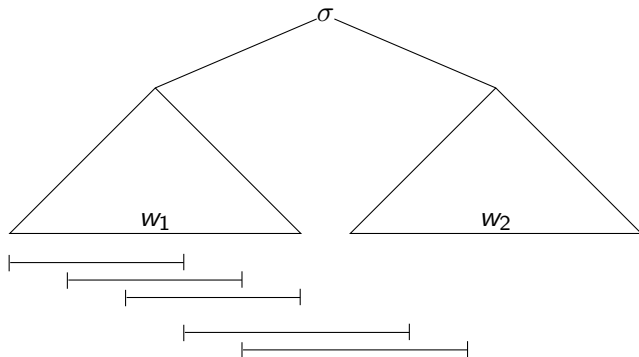
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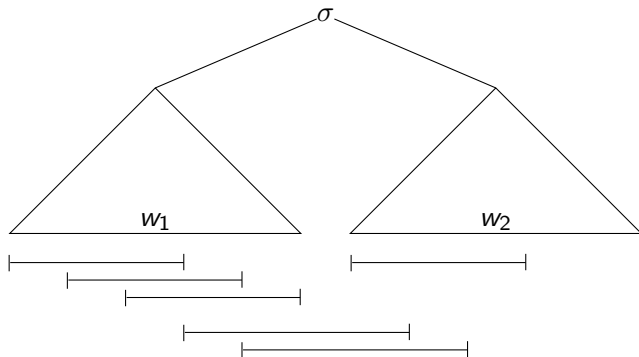
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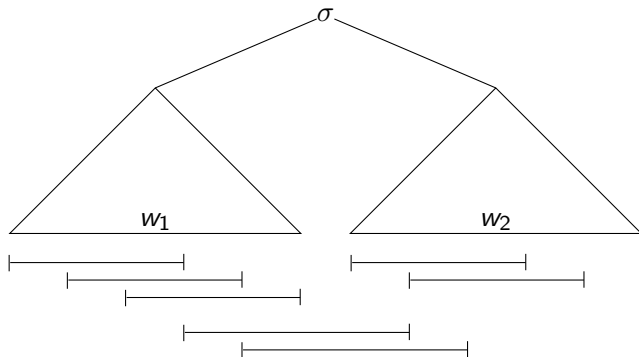
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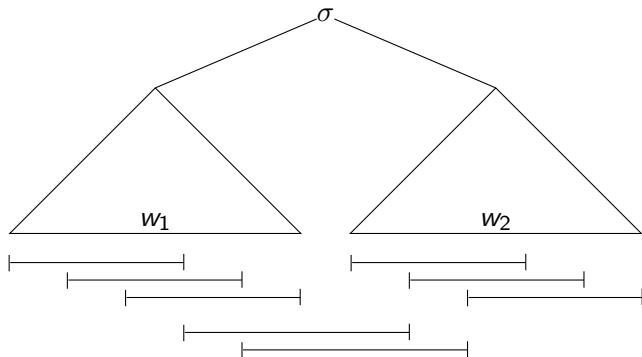
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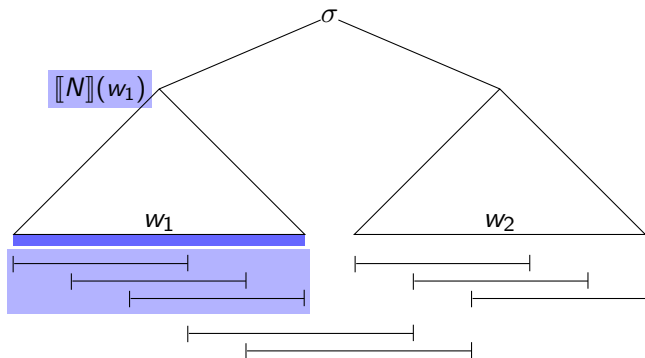
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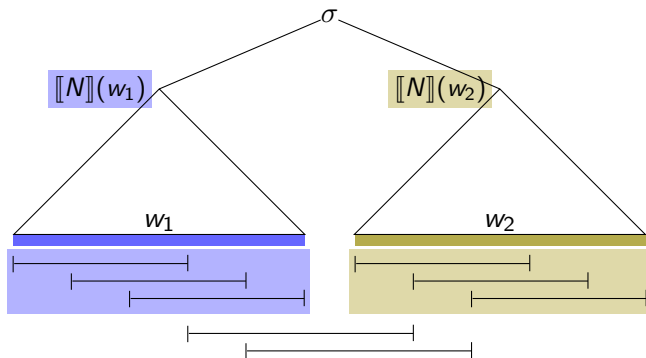


$$[[N]](w_1 w_2) = \boxed{[[N]](w_1)} \cdot \quad \cdot$$

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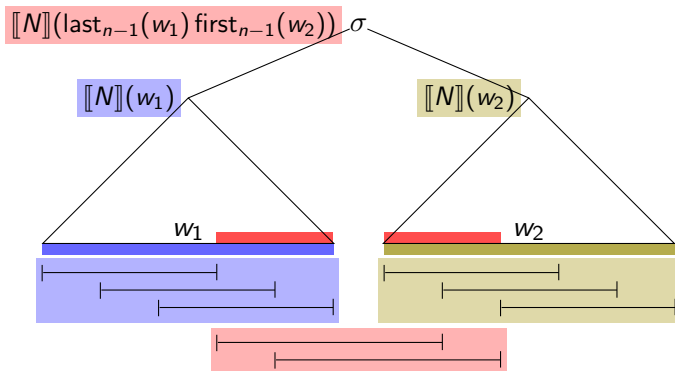
$$\llbracket N \rrbracket(w_1 w_2) = \llbracket N \rrbracket(w_1) \cdot \llbracket N \rrbracket(w_2)$$

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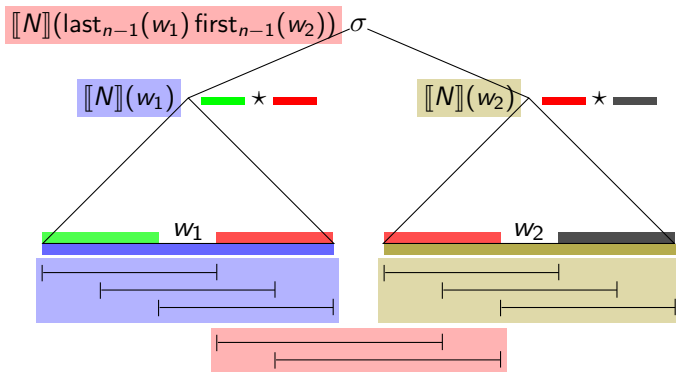


$$\llbracket N \rrbracket(w_1 w_2) = \llbracket N \rrbracket(w_1) \cdot \llbracket N \rrbracket(\text{last}_{n-1}(w_1) \text{ first}_{n-1}(w_2)) \cdot \llbracket N \rrbracket(w_2)$$

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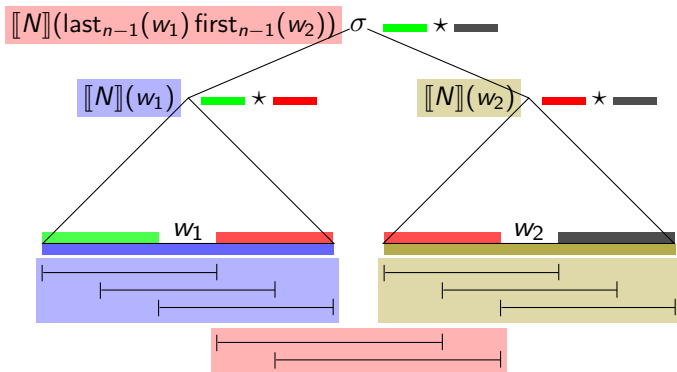


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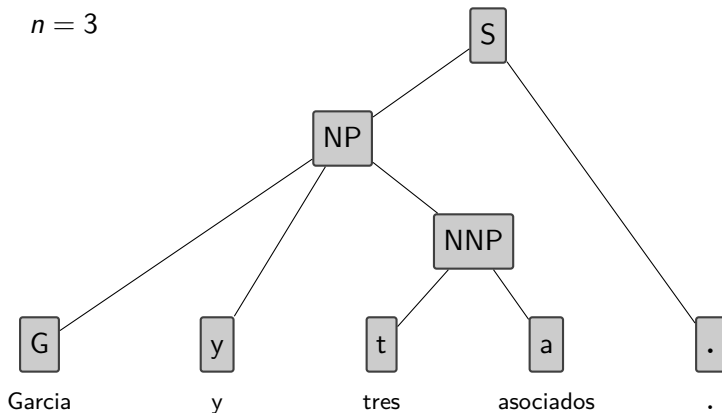


$$[[N]](w_1 w_2) = [[N]](w_1) \cdot [[N]](\text{last}_{n-1}(w_1) \text{ first}_{n-1}(w_2)) \cdot [[N]](w_2)$$

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# Generalized $n$ -Gram WTA (Example)

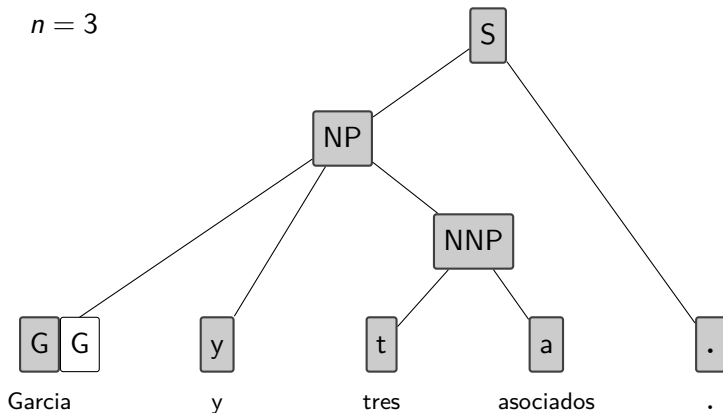
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$wt(\kappa) =$

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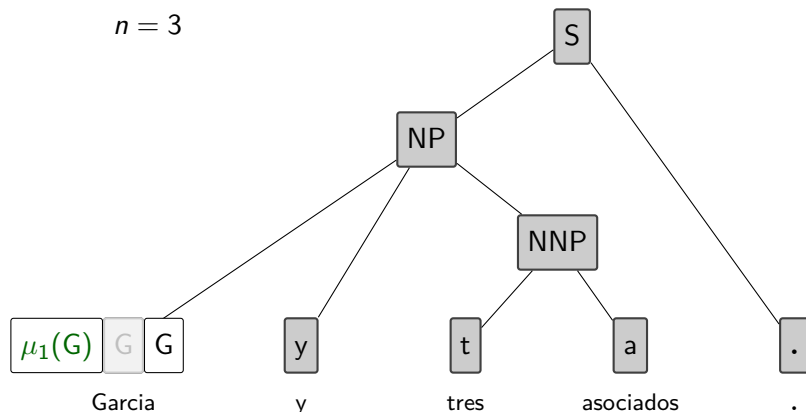
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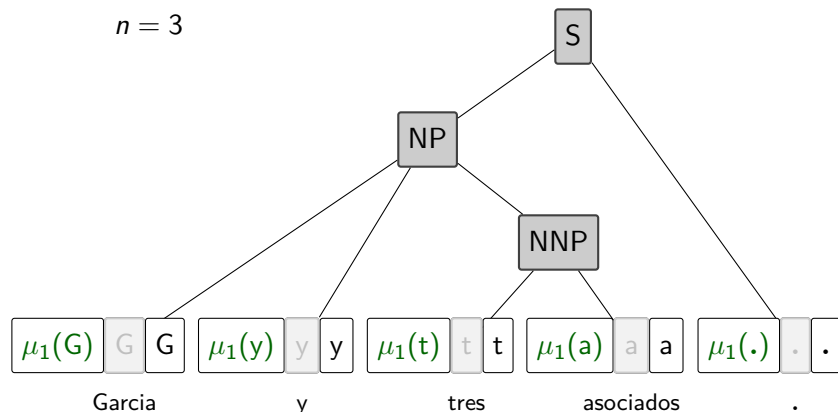
$n = 3$



$$wt(\kappa) = \mu_1(G)$$

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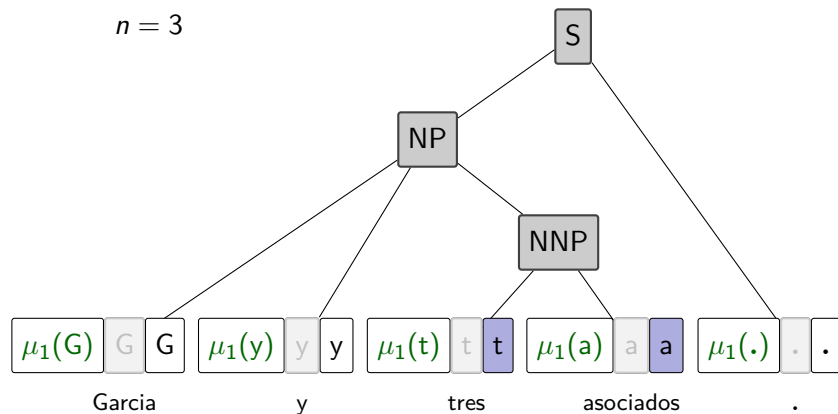
$n = 3$



$$\text{wt}(\kappa) = \mu_1(\text{G}) \cdot \mu_1(\text{y}) \cdot \mu_1(\text{t}) \cdot \mu_1(\text{a}) \cdot \mu_1(\text{.})$$

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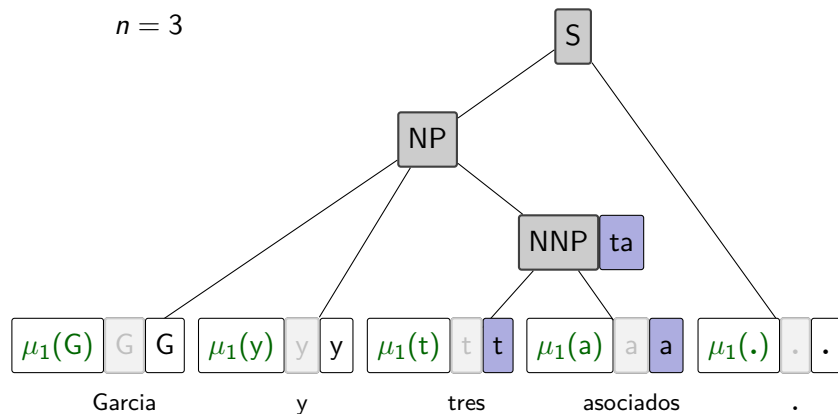


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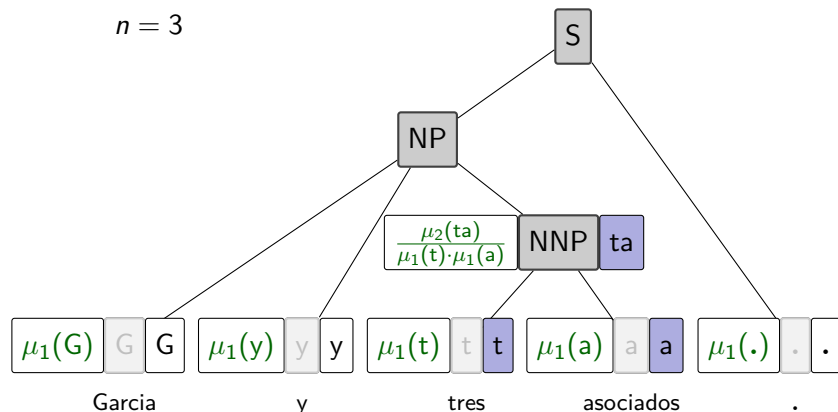
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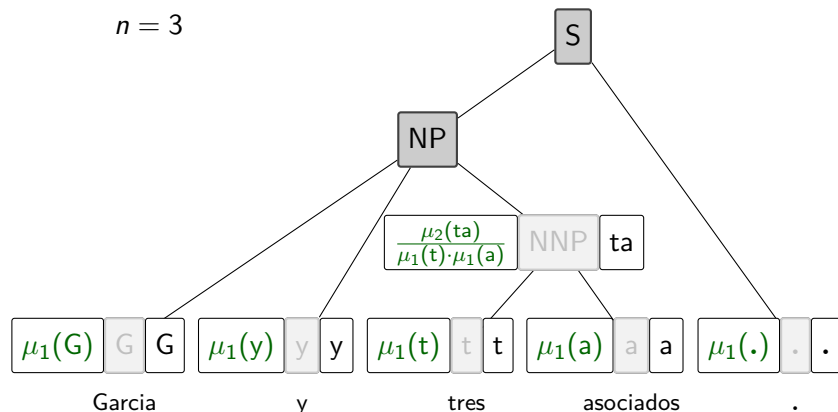
$n = 3$



$$wt(\kappa) = \mu_1(G) \cdot \mu_1(y) \cdot \mu_1(t) \cdot \mu_1(a) \cdot \mu_1(\cdot) \cdot \frac{\mu_2(ta)}{\mu_1(t) \cdot \mu_1(a)}$$

# Generalized $n$ -Gram WTA (Example)

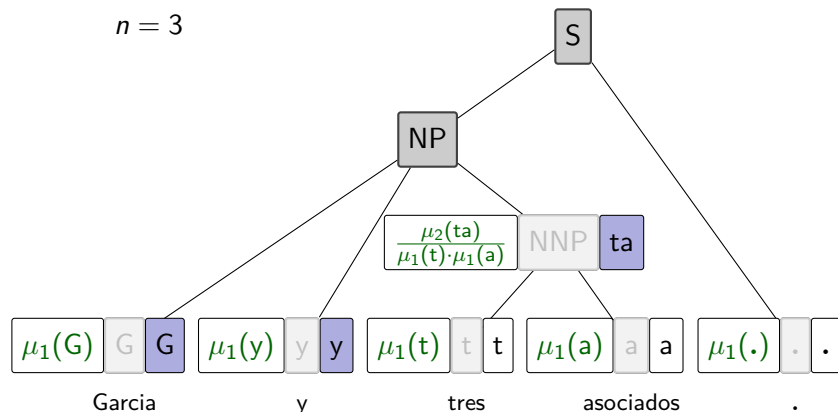
$n = 3$



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# Generalized $n$ -Gram WTA (Example)

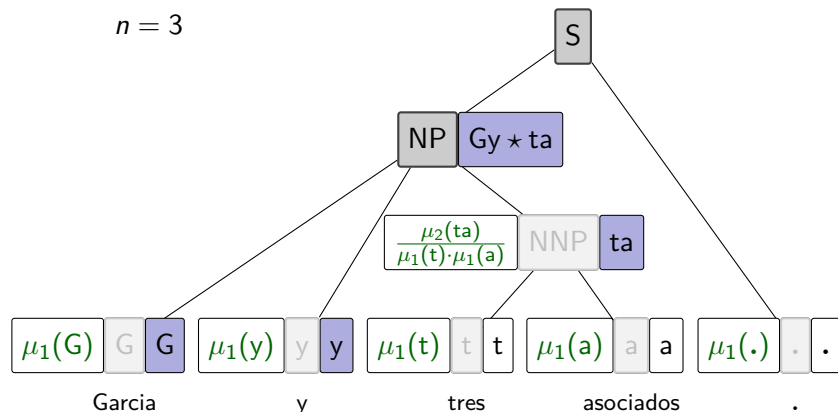
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$$wt(\kappa) = \mu_1(\text{G}) \cdot \mu_1(\text{y}) \cdot \cancel{\mu_1(\text{t})} \cdot \cancel{\mu_1(\text{a})} \cdot \mu_1(\text{.}) \cdot \frac{\mu_2(\text{ta})}{\mu_1(\text{t}) \cdot \mu_1(\text{a})}$$

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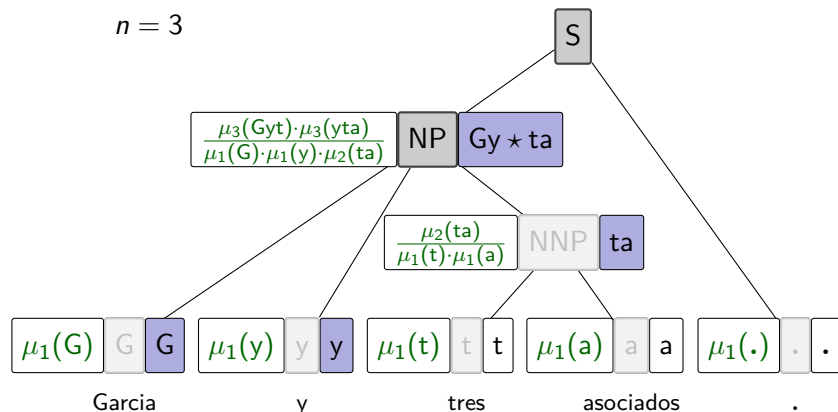
$n = 3$



$$wt(\kappa) = \mu_1(G) \cdot \mu_1(y) \cdot \cancel{\mu_1(t)} \cdot \cancel{\mu_1(a)} \cdot \mu_1(.) \cdot \frac{\mu_2(ta)}{\mu_1(t) \cdot \mu_1(a)}$$

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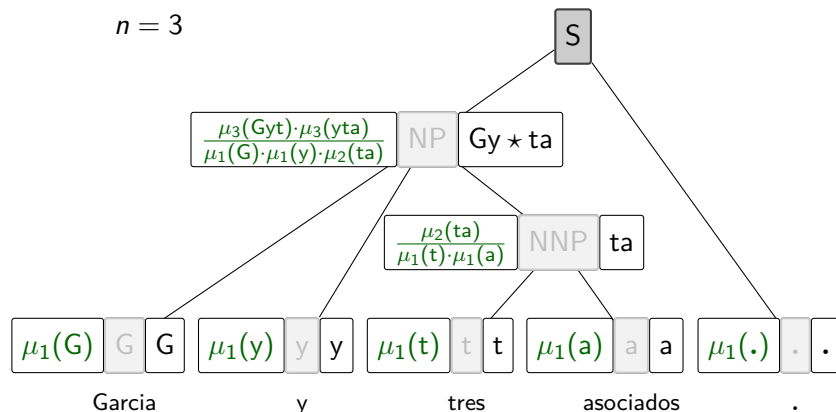
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# Generalized $n$ -Gram WTA (Example)

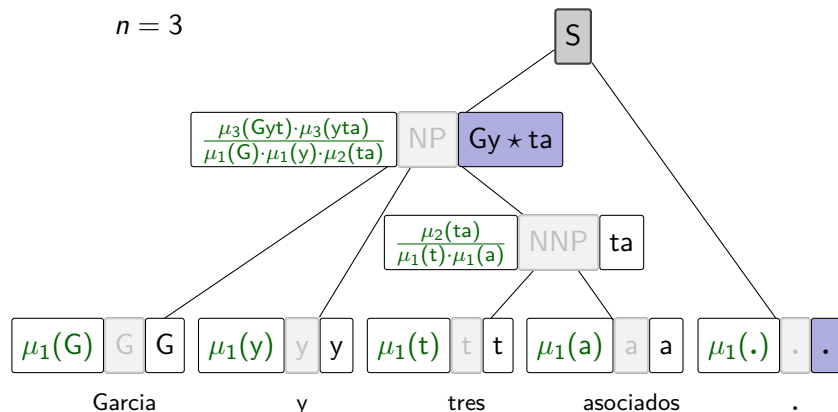
$n = 3$



$$wt(\kappa) = \cancel{\mu_1(G)} \cdot \cancel{\mu_1(y)} \cdot \cancel{\mu_1(t)} \cdot \cancel{\mu_1(a)} \cdot \mu_1(\cdot) \cdot \frac{\mu_2(ta)}{\mu_1(t) \cdot \mu_1(a)} \cdot \frac{\mu_3(Gyt) \cdot \mu_3(yta)}{\mu_1(G) \cdot \mu_1(y) \cdot \mu_2(ta)}$$

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$n = 3$

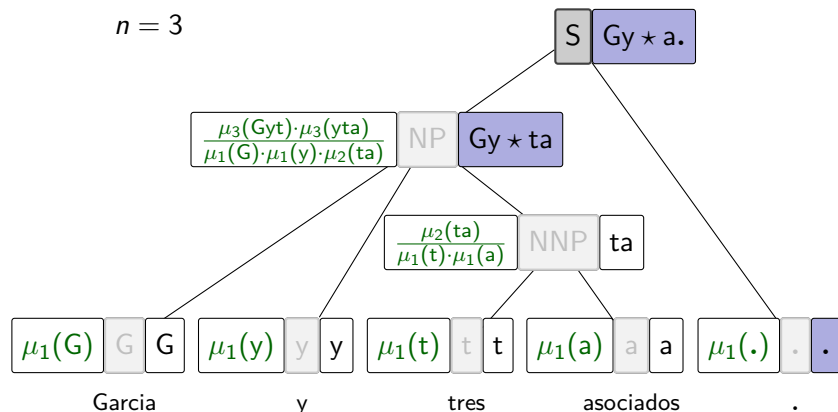


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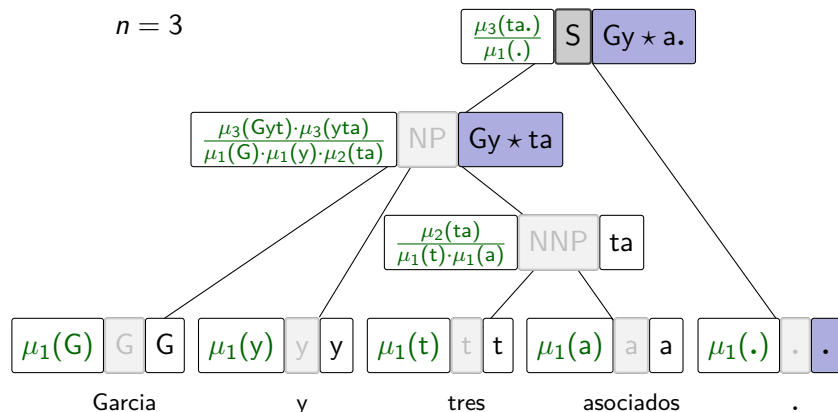
$n = 3$



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$n = 3$

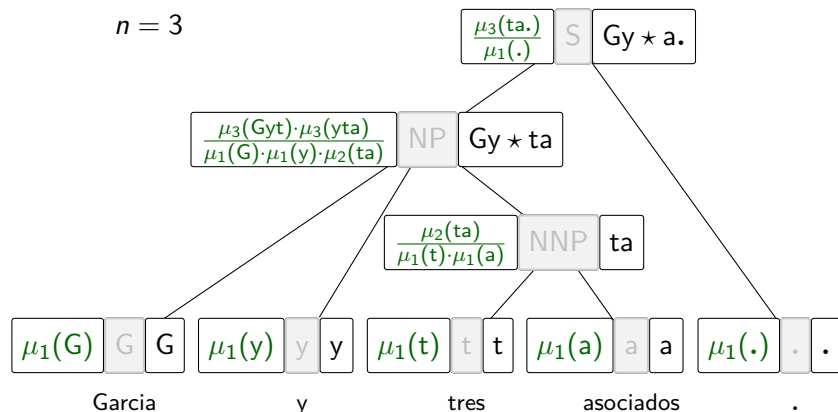


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# Generalized $n$ -Gram WTA (Example)

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$$\begin{aligned}
 wt(\kappa) &= \cancel{\mu_1(G)} \cdot \cancel{\mu_1(y)} \cdot \cancel{\mu_1(t)} \cdot \cancel{\mu_1(a)} \cdot \cancel{\mu_1(.)} \cdot \frac{\mu_2(ta)}{\mu_1(t) \cdot \mu_1(a)} \cdot \frac{\mu_3(Gyt) \cdot \mu_3(yta)}{\mu_1(G) \cdot \mu_1(y) \cdot \mu_2(ta)} \cdot \frac{\mu_3(ta.)}{\mu_1(.)} \\
 &= \mu_3(Gyt) \cdot \mu_3(yta) \cdot \mu_3(ta.) = \llbracket N \rrbracket(Gyta.)
 \end{aligned}$$

# Generalized $n$ -Gram WTA

The WTA  $\mathcal{A}_{N,\Sigma}$  is the tuple  $(Q, \Sigma, \delta, F)$  with

- $Q = \Gamma^{\leq n-1} \cup (\Gamma^{n-1} \times \{\star\} \times \Gamma^{n-1})$
- $F = Q$

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- $$\delta(q_1, \dots, q_k, \sigma, q) = \begin{cases} g(\sigma) & \text{if } k = 0 \text{ and } q = f(\sigma) \\ \frac{g(q_1 \dots q_k)}{g'(q_1) \dots g'(q_k)} & \text{if } k \geq 1 \text{ and } q = f(q_1 \dots q_k) \\ 0 & \text{otherwise} \end{cases}$$

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for every  $v \in (\Gamma \cup \{\star\})^*$

$$f(v) = \begin{cases} v & \text{if } v \in \Gamma^* \text{ and } |v| < n \\ v_1^{n-1} \star v_{|v|-n+2}^{|v|} & \text{if } (v \in \Gamma^* \text{ and } |v| \geq n) \text{ or} \\ & (v \notin \Gamma^* \text{ and } |v| \geq n) \\ \dots & \text{otherwise} \end{cases}$$

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for every  $v \in (\Gamma \cup \{\star\})^*$ ,  $v = u_0 \star u_1 \dots \star u_l$ ,  $u_i \in \Gamma^*$  for all  $i \in \{0, \dots, l\}$ ,  $w \in \Sigma^*$ , and  $q \in Q$

$$g(v) = \begin{cases} \llbracket N \rrbracket(u_0) & \text{if } l = 0 \\ N'(u_0) \cdot (\prod_{i=1}^{l-1} \llbracket N \rrbracket(u_i)) \cdot N'(u_l) & \text{otherwise} \end{cases}$$

$$N'(w) = \begin{cases} \llbracket N \rrbracket(w) & \text{if } |w| \geq n \\ 1 & \text{otherwise} \end{cases} \quad g'(q) = \begin{cases} \llbracket N \rrbracket(q) & \text{if } q \in \Gamma^* \\ 1 & \text{otherwise} \end{cases}$$



## Theorem

Let  $\Sigma$  be a ranked alphabet and  $N$  a generalized  $n$ -gram model over  $\Gamma$ .  
Then

- $\llbracket \mathcal{A}_{N,\Sigma} \rrbracket = \text{yield}^{-1}(\llbracket N \rrbracket)$
- $\mathcal{A}_{N,\Sigma}$  is bottom-up deterministic

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## Conjecture

$\mathcal{A}_{N,\Sigma}$  is well suited for the calculation of  $\text{argmax}_d ((f \triangleleft \text{TM}) \triangleright \llbracket N \rrbracket)(d)$ .

- Implement a decoder for SMT based on
  - BHPS, and
  - the generalized  $n$ -gram WTA.

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- Compare the runtime behaviour in practice.

- Bar-Hillel, Yehoshua, Micha Perles, and Eliahu Shamir (1961). “On Formal Properties of Simple Phrase Structure Grammars”. In: *Z. Phonetik. Sprach. Komm.* 14, pp. 143–172.
- Chiang, David (2007). “Hierarchical phrase-based translation”. In: *Computational Linguistics* 33(2), pp. 201–228.
- Maletti, Andreas and Giorgio Satta (2009). “Parsing Algorithms Based on Tree Automata”. In: *Proc. of IWPT '09*. Paris, France: ACL, pp. 1–12.

# BHSP for Trees and $n$ -Gram Models

$\mathcal{L} = T_\Sigma$  as **WTA**

$Q_1 = \{\star\}$

for every  $k \in \mathbb{N}, \sigma \in \Sigma^{(k)}$ :

$\star \xrightarrow{1} \sigma(\underbrace{\star, \dots, \star}_{k \text{ times}})$

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$R = \llbracket (\Gamma, \mu) \rrbracket$  as **WFSA**

$$Q_2 = \Gamma^0 \cup \dots \cup \Gamma^{n-1}$$

for every  $q \in Q \setminus \Gamma^{n-1}, \sigma \in \Gamma$ :

$$q \xrightarrow{\sigma/1} q\sigma$$

for every  $q \in \Gamma^{n-2}, \sigma, \sigma' \in \Gamma$ :

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$\mathcal{L} \odot R$  as **WTA**

$$Q = Q_2 \times Q_1 \times Q_2$$

for every  $q \in Q_2 \setminus \Gamma^{n-1}, \sigma, \sigma' \in \Gamma$ :

$$q \star q\sigma \xrightarrow{1} \sigma, \quad \sigma'q \star q\sigma \xrightarrow{\mu(\sigma'q\sigma)} \sigma$$

for every  $k \in \mathbb{N}, \sigma \in \Sigma^{(k)} \setminus \Gamma, q_0, \dots, q_k \in Q_2$ :

$$q_0 \star q_k \xrightarrow{1} \sigma(q_0 \star q_1, q_1 \star q_2, \dots, q_{k-1} \star q_k)$$

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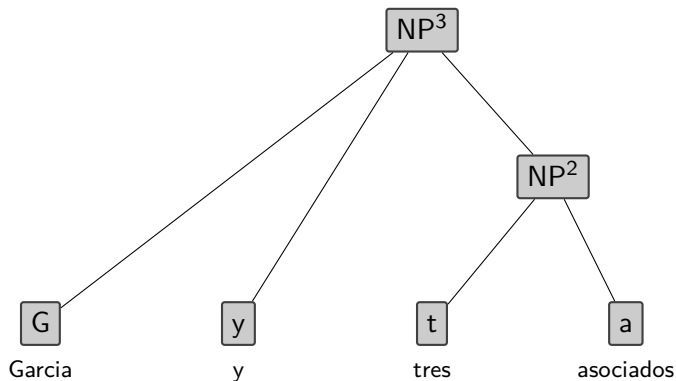
**WFSA**. . . weighted finite state automaton

$$\Gamma = \Sigma^{(0)}$$



# BHSP for Trees and $n$ -Gram Models (Example)

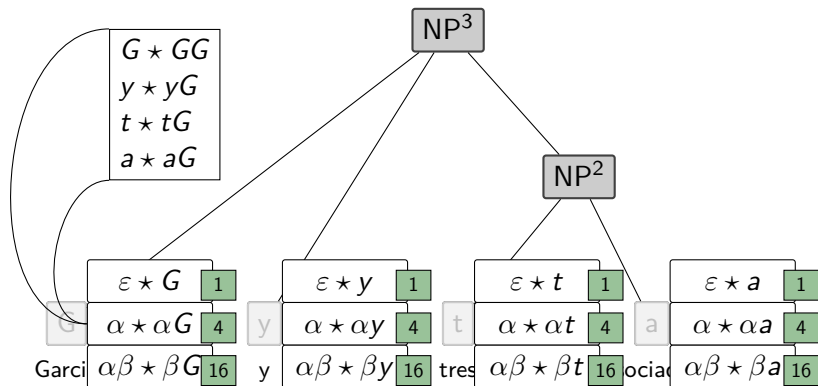
$$n = 3$$
$$\Gamma = \{G, y, t, a\}$$



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$$n = 3$$

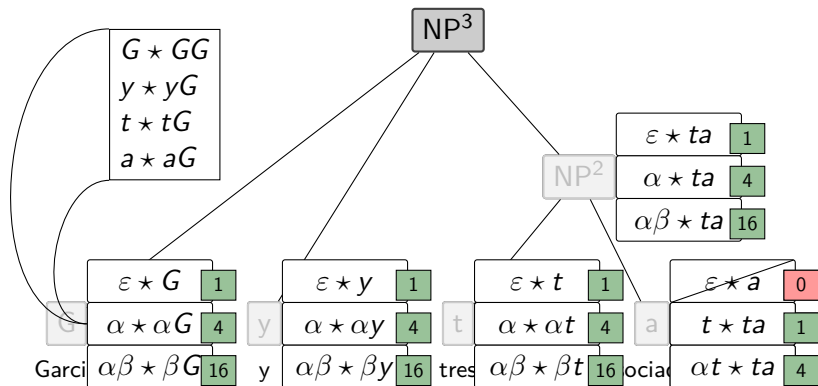
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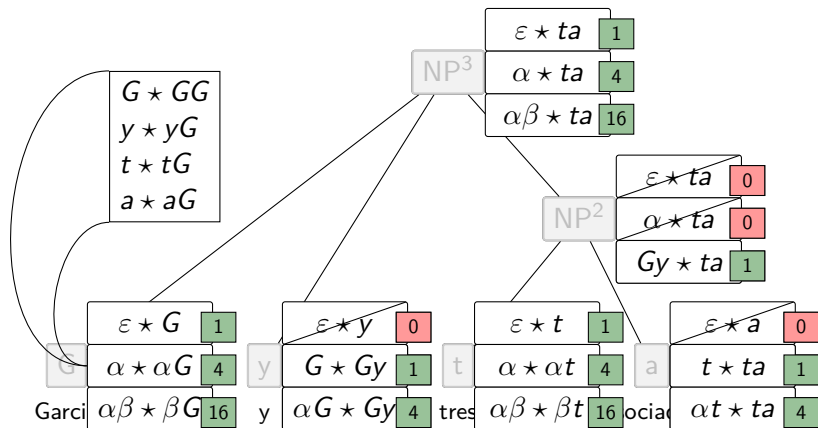
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