

# Group testing with geometric ranges

Benjamin Aram Berendsohn



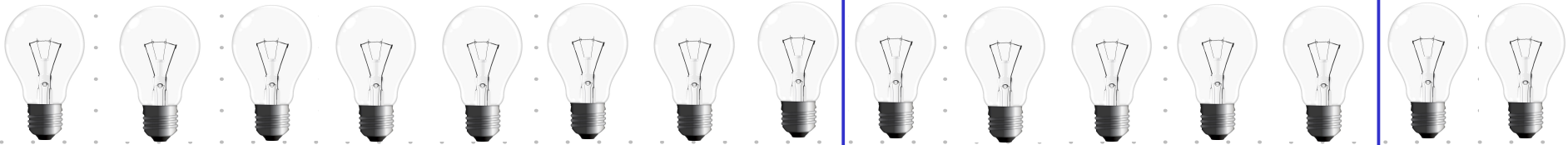
László Kozma



Freie Universität Berlin

ISIT 2022

# Group testing



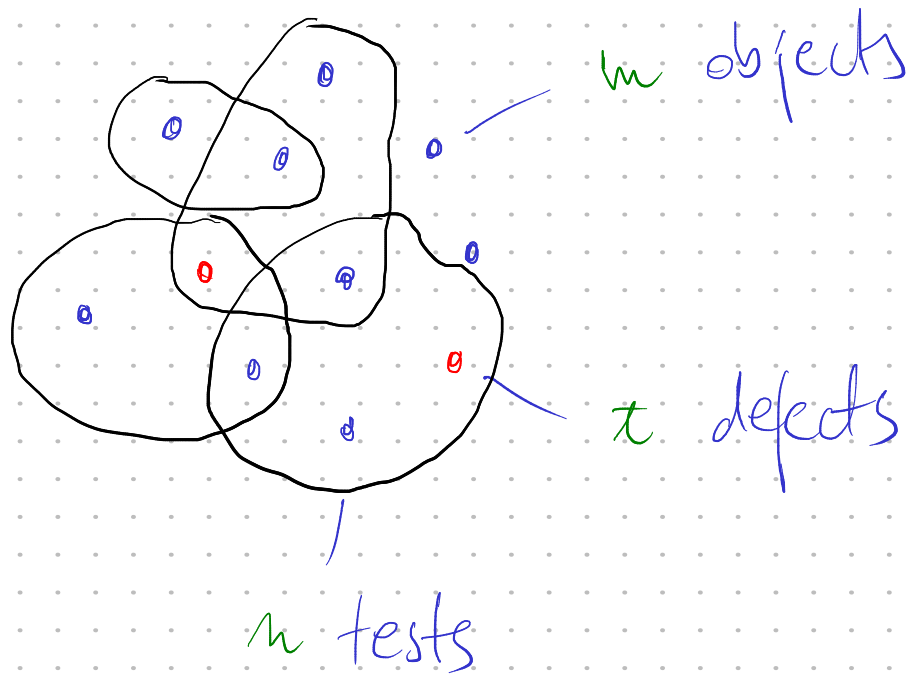
test: are these **all** good?  
YES/NO

$m$  objects

$n$  tests (non-adaptive)

$t$  defects

→ Can a given configuration  
of objects/tests identify  
 $t$  defective objects?



→ non-adaptive test configuration:  
 can be viewed as a  
 set system

→ Can a given configuration  
 identify  $t$  defectives?

→ For given  $t, n$  what is largest  $m$  we can handle?

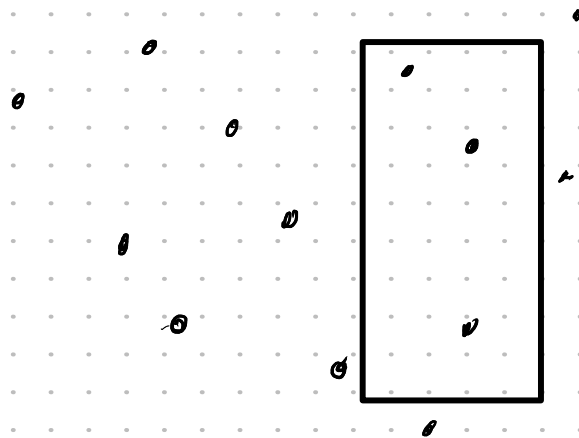
Classical group testing: for  $t \in O(1)$ ,  $m = 2^{\Theta(n)}$

(tests unrestricted)

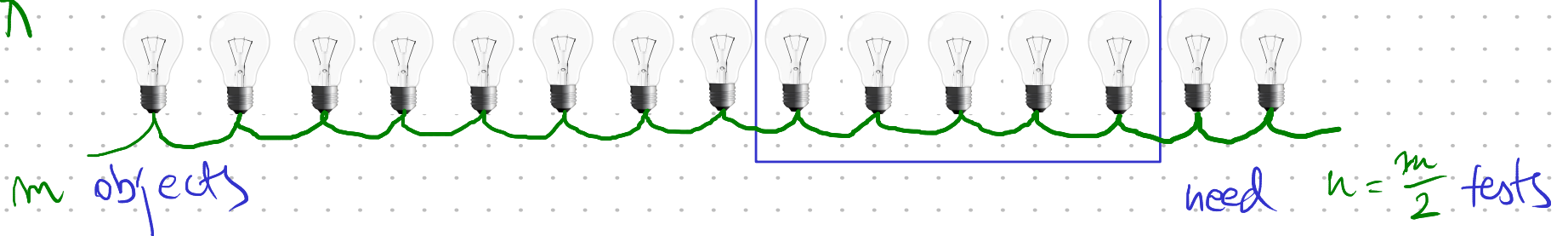
Our setting:

→ Objects are points in  $\mathbb{R}^d$

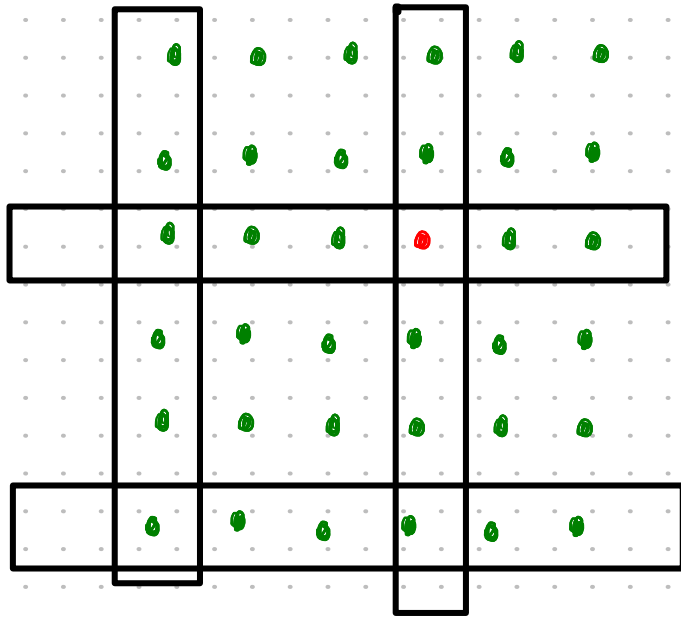
→ Tests are rectangles in  $\mathbb{R}^d$  (parameter ranges)



e.g.  $\mathbb{R}^1$



e.g.



$m$  objects

$t = 1$  defects

$n = 2\sqrt{m}$  tests

→ How many objects (points) can we handle  
with  $n$  rectangle-tests,  $t$  defects, in  $\mathbb{R}^d$ ?

$S_t^d(n)$ : max # of points such that there  
exists a valid configuration

$$S_t^d(n)$$

Valid configuration:

Def. for every two sets  $A, B$  such that  $|A| = |B| = t$

$\{\text{set of tests intersecting } A\} \neq \{\text{set of tests intersecting } B\}$

(can distinguish cases "A is defect set" and "B is defect set")

$\rightarrow t$ -separable

$S_t^d(u)$  → goal: understand asymptotic behavior.

Note: does not imply an efficient algorithm for identifying defects.

for that we have a stronger property →  $t$ -disjunct

Similar quantity: →  $D_t^d(u)$

Obs.  $D_t^d(u) \leq S_t^d(u)$

Goal: asymptotics of  $S_t^d(n)$  (max # of points with  
 $n$  rectangle-tests  
 $t$  defects  
 $d$  dimensions)

Recall combinatorial setting:  $2^{\Theta(n)}$

Easy observations:

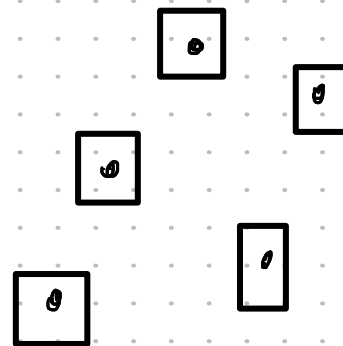
$$S_t^d(n) \geq S_{t+1}^d(n)$$

$$S_t^d(n) \leq S_t^{d+1}(n)$$

Trivial lower bound:

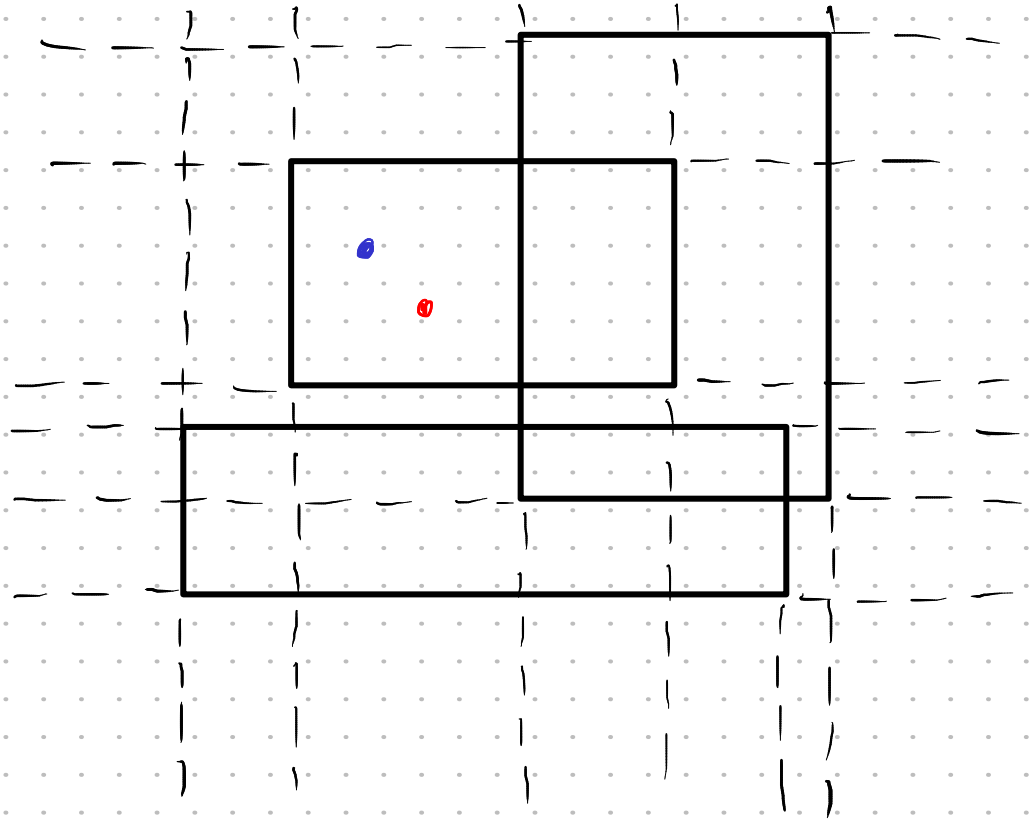
$$n \leq S_t^d(w)$$

→  $w$  tests can always handle  
at least  $w$  points:



(Almost) Trivial upper bound:

$$S_t^d(n) \leq n^d$$



→  $n$  rectangles define  $(2n)^d$  cells

→ points in same cell cannot be distinguished

$$n \leq S_t^d(u) \leq n^d$$

→ polynomial instead of exponential

→ What is the exponent?

Claim:  $S_1^d \sim n^d$

$$S_t^1 \sim n$$

Main result: neither LB nor UB tight for  $t, d \geq 2$ .

Result:

$$n^{1+f(t,d)} \leq S_t^d(n) \leq n^{d-g(t,d)}$$

→ nontrivial improvements  $f(\cdot), g(\cdot)$

→ gap grows with  $d, t$

eg.  $n^{3/2} \leq S_2^2 \leq n^{5/3}$

→ few tight results

eg.  $D_2^2 \sim n^{3/2}$

Open: close gaps, find separations.

Lower bound sketch:

$$n^{1+2/t} \leq S_t^2(n)$$

→ need to show a concrete feasible configuration

Idea:

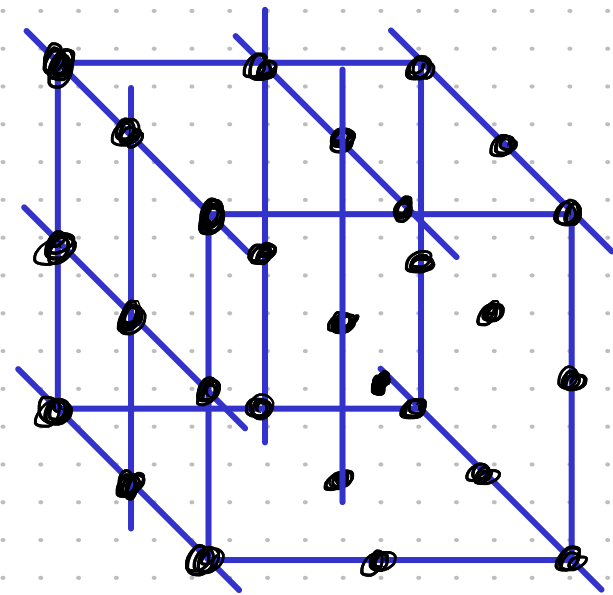
start with an abstract construction:

→  $k^d$  grid points

→  $d \cdot k^{d-1}$  lines as tests

→ exact characterization:

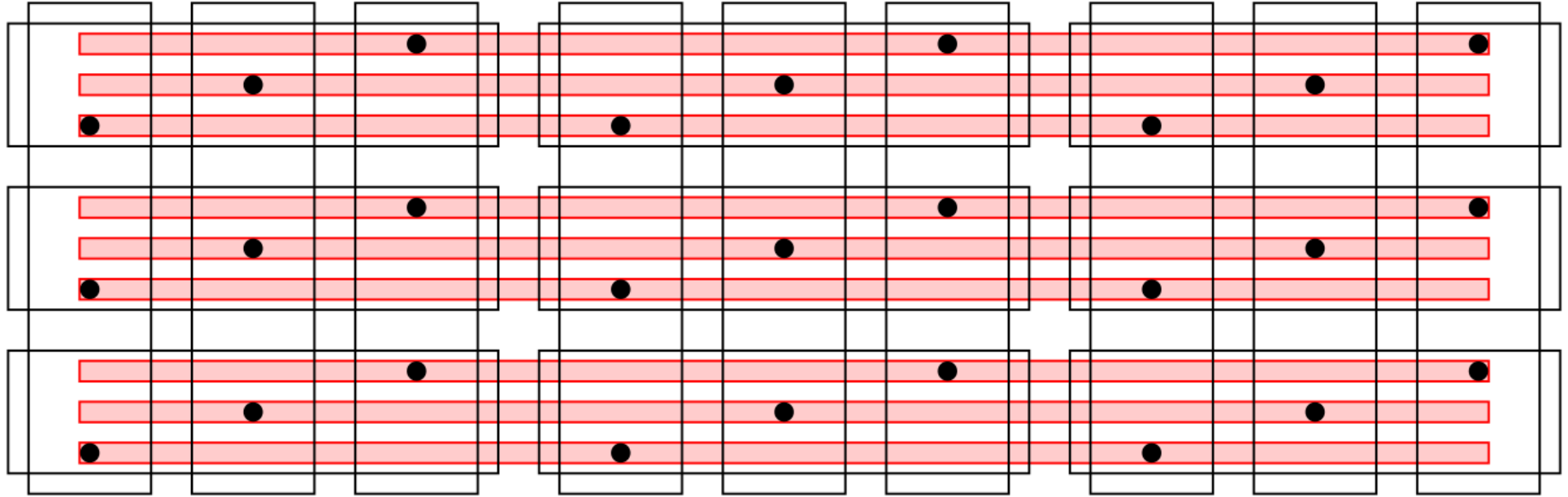
$\sim 2d$ -separable



→ Can "project" to equivalent rectangle-system in  $\mathbb{R}^2$

→ can set  $k$  as needed

e.g. 3-separable system in  $\mathbb{R}^2$

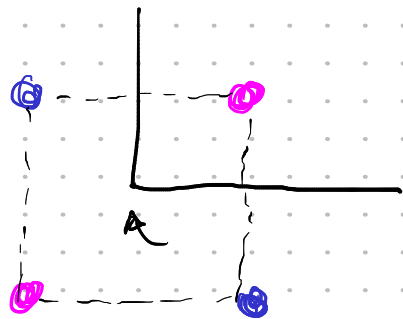


## Upper bounds

e.g.  $S_2^2(n) \in O(n^{5/3})$  (no valid configuration with more points)

→ discretize to  $4n \times 4n$  grid

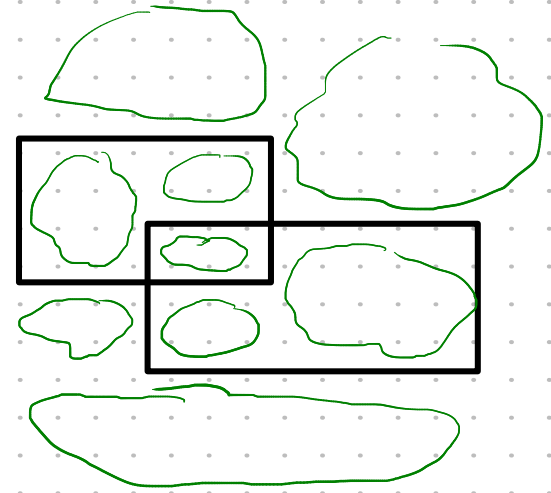
→ Observe: if 4 points (objects) form a rectangle, must have some test corner inside



otherwise cannot distinguish diagonal pairs.

→ Observe: if 4 points (objects) form a rectangle,  
must have some test corner inside

Proof idea: → partition space into regions  
without rectangle corners



⇒ these regions cannot have 4 points  
forming rectangles



⇒ these regions cannot have  
too many points

⇒ overall we cannot have  
too many points



Thm. [Kövari-Sós-Turán]  
Any  $n \times n$  matrix with  
at least  $n^{3/2}$  \*-entries

must contain:

*	—	*
*	—	*

Somewhere.

→ carefully control parameters

→ similar bounds by lifting other theorems from extremal combinatorics.

## Conclusion

- possibility and limitations of group testing with rectangles
- $n$  tests can handle  $n^c$  objects (as opposed to  $c^n$ )
- Lower bounds: concrete geometric constructions
- Upper bounds: impossibility via extremal combinatorics

Open:

- exact bounds elusive, large gaps remain
- asymptotic separations for different  $t, d$ ?

THANKS!